

MULTIMEDIA



UNIVERSITY

STUDENT ID NO

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# MULTIMEDIA UNIVERSITY

## FINAL EXAMINATION

TRIMESTER 2, 2015/2016

### TCT 2561 – COMPLEXITY THEORY

( All sections / Groups )

9 MARCH 2016

9:00 a.m. – 11:00 a.m.

( 2 Hours )

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#### INSTRUCTIONS TO STUDENTS

1. This Question paper consists of **5 pages** only including the cover page.
2. Attempt **ALL** questions.
3. All questions carry equal marks and the distribution of the marks for each question is given.
4. Please print all your answers **CLEARLY** in the Answer Booklet provided.

**Question 1** (2+2+6+5 marks)

- (a) Differentiate between computational complexity theory and computability theory.
- (b) Differentiate between deterministic Turing machine and non-deterministic Turing machine.
- (c) Draw a graph to illustrate each of the following asymptotic equations.
- i.  $f(n) = O(g(n))$
  - ii.  $f(n) = \Omega(g(n))$
  - iii.  $f(n) = \Theta(g(n))$
- (d) Consider the following algorithm for the Tower of Hanoi problem with  $n$  discs and three poles: `src` is the source, `spare` is the spare, and `dest` is the destination.

```
void towerOfHanoi(int n, char src, char spare, char dest)
{
    towerOfHanoi (n - 1, src, dest, spare);
    towerOfHanoi (1, src, spare, dest);
    towerOfHanoi (n - 1, spare, src, dest);
}
```

- i. Write down the recurrence relation of `towerOfHanoi ()`.
- ii. What is the time complexity class of `towerOfHanoi ()` and why?

Continued .....

**Question 2** (3+8+4 marks)

- (a) Savitch's theorem says that any non-deterministic Turing machine that uses  $f(n)$  space can be converted to a deterministic Turing machine that uses only  $f^2(n)$  space.
- Give the formal definition of Savitch's theorem.
  - What is the significance of Savitch's theorem?
  - What is the implication on the time complexity?
- (b) Given the following definition.

$PATH = \{ \langle G, s, t \rangle \mid G \text{ is a directed graph that has a directed path from } s \text{ to } t \}.$

- Give a high-level Turing machine description for  $PATH$  in exponential time.
  - Give a high-level Turing machine description for  $PATH$  in polynomial time.
  - Is  $PATH$  problem in time complexity class P? Why?
- (c) Given the following definitions.

$TQBF = \{ \langle \phi \rangle \mid \phi \text{ is a true fully quantified Boolean formula} \}.$

$FORMULA-GAME = \{ \langle \phi \rangle \mid \text{Player } E \text{ has a winning strategy in the formula game associated with } \phi \}.$

Show that  $FORMULA-GAME$  is PSPACE-complete.

Continued .....

**Question 3** (8+3+4 marks)

- (a) Give a complexity class example for each of the following computational models and then briefly explain your reason.
- Boolean circuit
  - Probabilistic Turing machine
  - Alternation
  - Interactive proof system
- (b) Is  $\text{NPSPACE} \subseteq \text{TIME}(2^{n^k})$ ? Explain your decision.
- (c) Draw a Venn diagram that depicts the relationship between NP-complete, NPSPACE, NP, PSPACE, and NP-hard complexity classes. Label the complexity classes clearly in your drawing.

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**Question 4** (3+3+5+4 marks)

- (a) In your own words, explain the theorem, “if  $A \leq_P B$  and  $B \in P$  then  $A \in P$ ”.
- (b) Draw a figure to illustrate mapping reducibility.
- (c) Examine the following definitions.

$SORTING = \{ \langle A[], n \rangle \mid A[] \text{ is an array of integers and } n \text{ is the array size such that we have the array of integers in ascending order } \}.$

$DISTINCT = \{ \langle A[], n \rangle \mid A[] \text{ is an array of integers and } n \text{ is the array size such that we have distinct integers in the array } \}.$

Construct a polynomial time reduction from *SORTING* to *DISTINCT*.

- (d) Describe two methods to prove that a problem B is NP-complete.

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